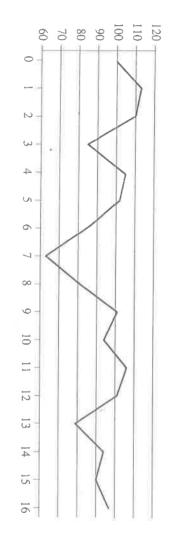
老 興 試科目 1-1-Algorithms

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Suppose that you been offered the opportunity to invest in the Volatile Chemical possible price and later on sell at the highest possible priceshare. Of course, you would want to "buy low, sell high"-buy at the lowest buy the stock at any one time, starting after day 0, when the price is \$100 per your profit. Figure 1 shows the price of the stock over a 17-day period. You may learn what the price of the stock will be in the future. Your goal is to maximize close of trading for the day. To compensate for this restriction, you are allowed to of stock only one time and then sell it at a later date, buying and selling after the occurs after day 7, which occurs after the highest price, after day 1. sell at the highest price within a given period. In Figure 1, the lowest price profit. Unfortunately, you might not be able to buy at the lowest price and then Volatile Chemical Corporation is rather volatile. You are allowed to buy one unit Corporation. Like the chemicals the company produces, the stock price of the to maximize your



Change	Price	Day
	100	C
13	113	-
3	011	1
-25	85	C.
20	105	4
5	102	Û
-16	86	6
-23	63	7
\propto	00	00
20	101	9
-7	94	10
12	106	11
5	101	12
-22	79	13
15	- 1	14
<u>-</u> 4	90	15
7	97	16

bottom row of the table gives the change in price from the previous day axis of the chart indicates the day, and the vertical axis shows the price. The Corporation after the close of trading over a period of 17 days. The horizontal Figure 1: Information about the price of stock in the Volatile Chemical

- (a) What is the major difference between "divide-and-conquer" and "dynamic programming"? (5%)
- **b** how to solve the problem in detail. (20%) Which technique should be applied to solve the problem? Please explain

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a search will correspond to d_i . Figure 2 shows a binary search trees for a set of nvalues between k_i and k_{i+1} . For each dummy key d_i , we have a probability q_i that values greater than k_n , and for i = 1, 2, ..., n - 1, the dummy key d_i represents all not in K. In particular, d_0 represents all values less than k_1 , d_n represents all In the *optimal binary search tree* problem, we are given a sequence $K = \langle k_1, k_2 \rangle$ K, and so we also have n+1 "dummy keys" $d_0, d_1, d_2, ..., d_n$ representing values probability p_i that a search will be for k_i . Some searches may be for values not in wish to build a binary search tree from these keys. For each key k_i , we have a ..., $k_n > \text{ of } n$ distinct keys in sorted order (so that $k_1 < k_2 < ...$ $< k_n$), and we

q_i	p_i	;	= 5
0.05		0	ceys wi
0.10	0.15	ш	keys with the following probabi
0.05	0.10	2	ollowing
0.05	0.05	ω	g probal
0.05	0.10	4	oilities:
0.10	0.20	S	

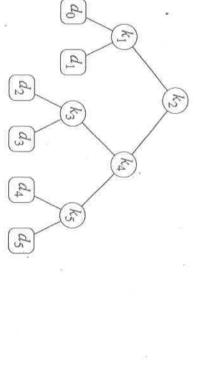


Figure 2: A binary search tree with expected search cost 2.80

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米 興 試科目 1-1-科技大學 Algorithms

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key, we can determine the expected cost of a search in a given binary search tree key d_i). Because we have probabilities of searches for each key and each dummy either successful (finding some key k_i) or unsuccessful (finding some dummy Each key k_i is an internal node, and each dummy key d_i is a leaf. Every search is

examined, i.e., the depth of the node found by the search in
$$T$$
, plus 1. Then the expected cost of a search in T is
$$\mathbb{E}\left[\operatorname{search cost in} T\right] = \sum_{i=1}^{n} (\operatorname{depth}_{T}(k_{i}) + 1) \cdot p_{i} + \sum_{i=0}^{n} (\operatorname{depth}_{T}(d_{i}) + 1) \cdot q_{i} \\ = 1 + \sum_{i=1}^{n} \operatorname{depth}_{T}(k_{i}) \cdot p_{i} + \sum_{i=0}^{n} \operatorname{depth}_{T}(d_{i}) \cdot q_{i} \ .$$

T. Let us assume that the actual cost of a search equals the number of nodes

the expected search cost node by node: where depth $_T$ denotes a node's depth in the tree T. In Figure 2, we can calculate

											G	
Total	d_5	d_4	d_3	d_2	d_1 .	d_0	k_5	k_4	k_3	k_2	k_1	node
	w	ω	ω	w	2	2	2	_	2	0	4	depth
	0.10	0.05	0.05	0.05	0.10	0.05	0.20	0.10	0.05	0.10	0.15	probability
2.80	0.40	0.20	0.20	0.20	0.30	0.15	0.60	0.20	0.15	0.10	0.30	contribution

search cost is smallest. We call such a tree an optimal binary search tree However, the tree in Figure 2 is not an optimal binary search tree. For the given set of probabilities, we wish to construct a binary search tree whose expected

- (a) What is the major difference between "dynamic programming" and "greedy algorithm"? (5%)
- 9 Which technique should be applied to solve the problem? Please explain how to solve the problem in detail. (20%)
- S Given a B tree with degree t,
- (a) What is the minimal number of children and keys for each internal node except the root node? (4%)
- (b) Given an ordered data sequence: 30, 2, 10, 20, 80, 60, 70, 92, 100, 55, 77, please construct a B tree with t=2 by inserting the data sequence. (6%)
- (c) After (b), please delete the keys 30, 55, 80, 10 sequentially from the B tree and draw the final result. (10%)

	(b) Based on the post-order sequence, please sort the numbers by insertion sort and merge sort, respectively. (You should write down all the steps) (10%)
	6. Given an ordered number sequence as follows: 50, 83, 10, 19, 5, 78, 39, 99, 77, 18, 1, 8. (a) Please construct a binary search tree by referring to the number sequence.
	a 1 13 h 7 i
	4. Asymptotic notations are usually be used to quantify the time complexity in a systematic way. Please make a comparison of the following big-O notations: $O(n^2)$, $O(\log n)$, $O(2^n)$, $O(n!)$, $O(n\log n)$, $O(n)$. (5%) 5. Given a connected weighted graph, please draw the minimal spanning trees by using Prim's and Kruskal's algorithms, respectively. (10%)
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